



# Anticipatory Routing For Highly Mobile Endpoints

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# Outline

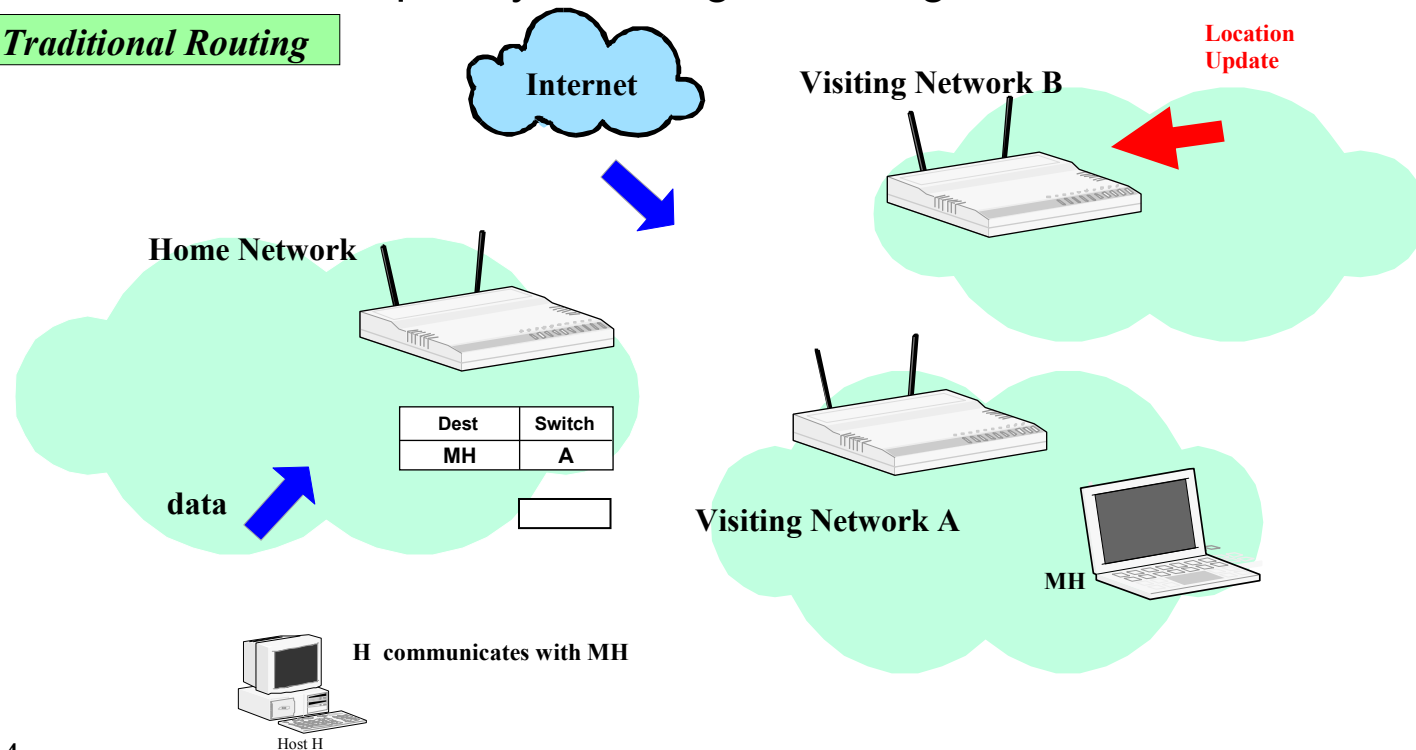
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# Problem Statement & Motivation

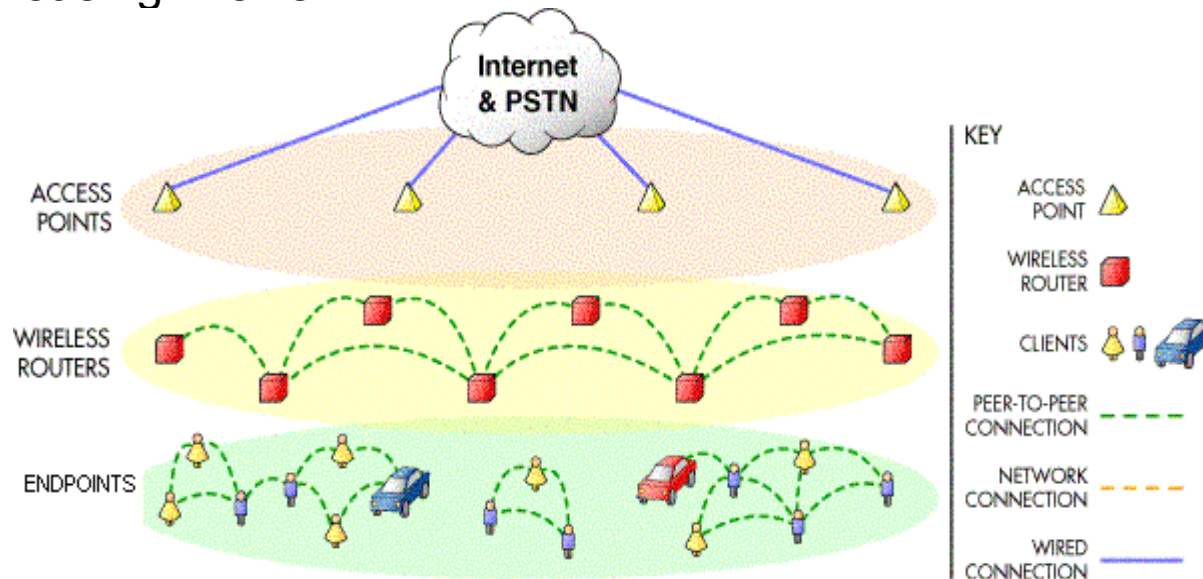
- All Mobile networks employ some control mechanism to route packets to mobile hosts.
  - Efficiency of control mechanism depends on two factors:
    - frequency of location change
    - time it takes to the system to know about that change.
  - But what if the frequency of change is so high?

## Traditional Routing



# Problem Statement & Motivation (cont'd)

- Wireless network architectures are evolving toward short links
  - Causes: need for high data rate, scarcity of communication spectrum, need to conserve battery, hence transmission power.
  - Consequence: increase in frequency of changes in network attachment points.
- Wireless network infrastructures (e.g wireless mesh networks/802.16) are increasing in size



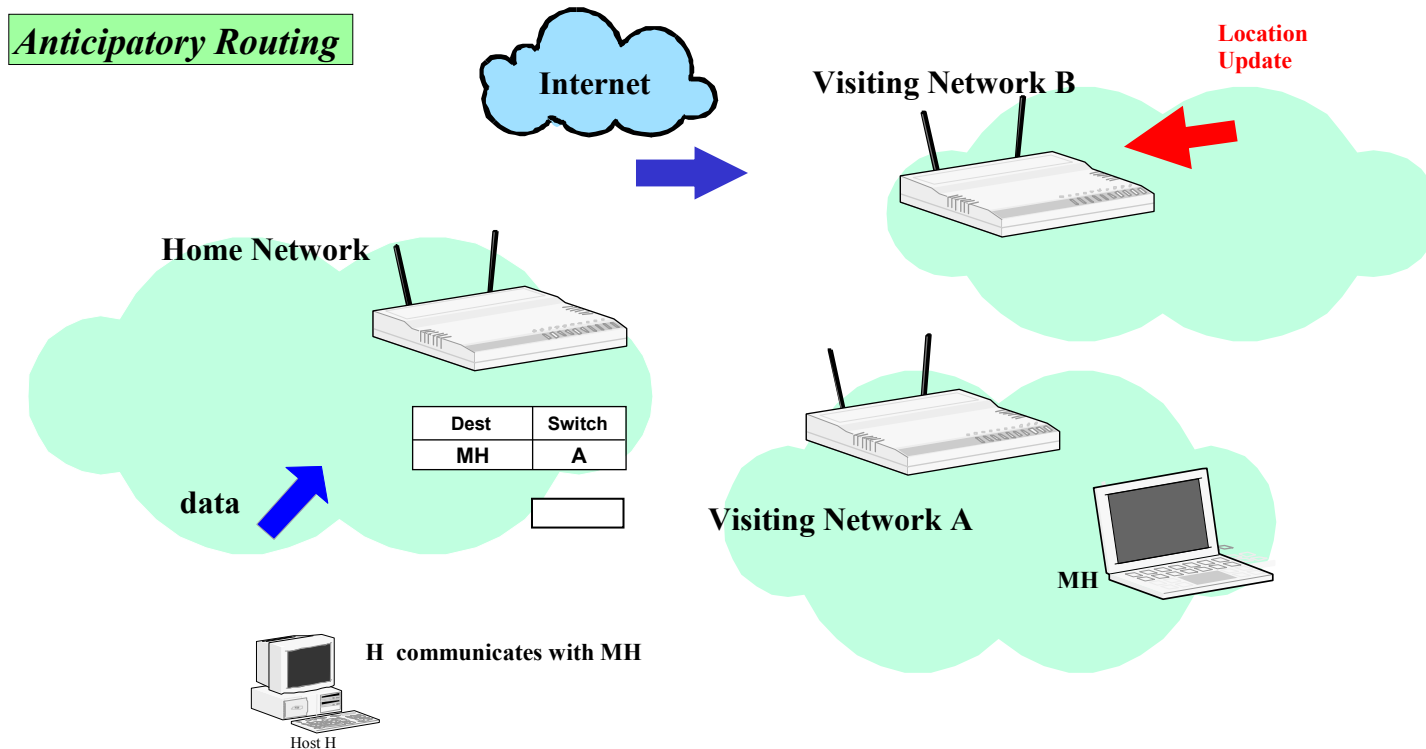
- Consequence: increase in control message latency, therefore increase in convergence time of the control mechanism.

## Problem Statement & Motivation (cont'd)

- Higher frequency of change and higher time of convergence may lead to situations where control mechanism is constantly outpaced by endpoint's mobility
- It is not the absolute speed that matters, but its relation to cell size and convergence time of the system, hence notion of *effective mobility*
- *Effective mobility* ( $M_{eff}$ ) = *convergence time / cell presence time*
  - When  $M_{eff} > 1$ , the system is constantly outpaced by endpoint's mobility and we say that the *reactive limit* has been reached

# Problem Statement & Motivation (cont'd)

- For mobility regimes that take system beyond reactive limit, a predictive (or anticipatory) mechanism is required.



# Related Work

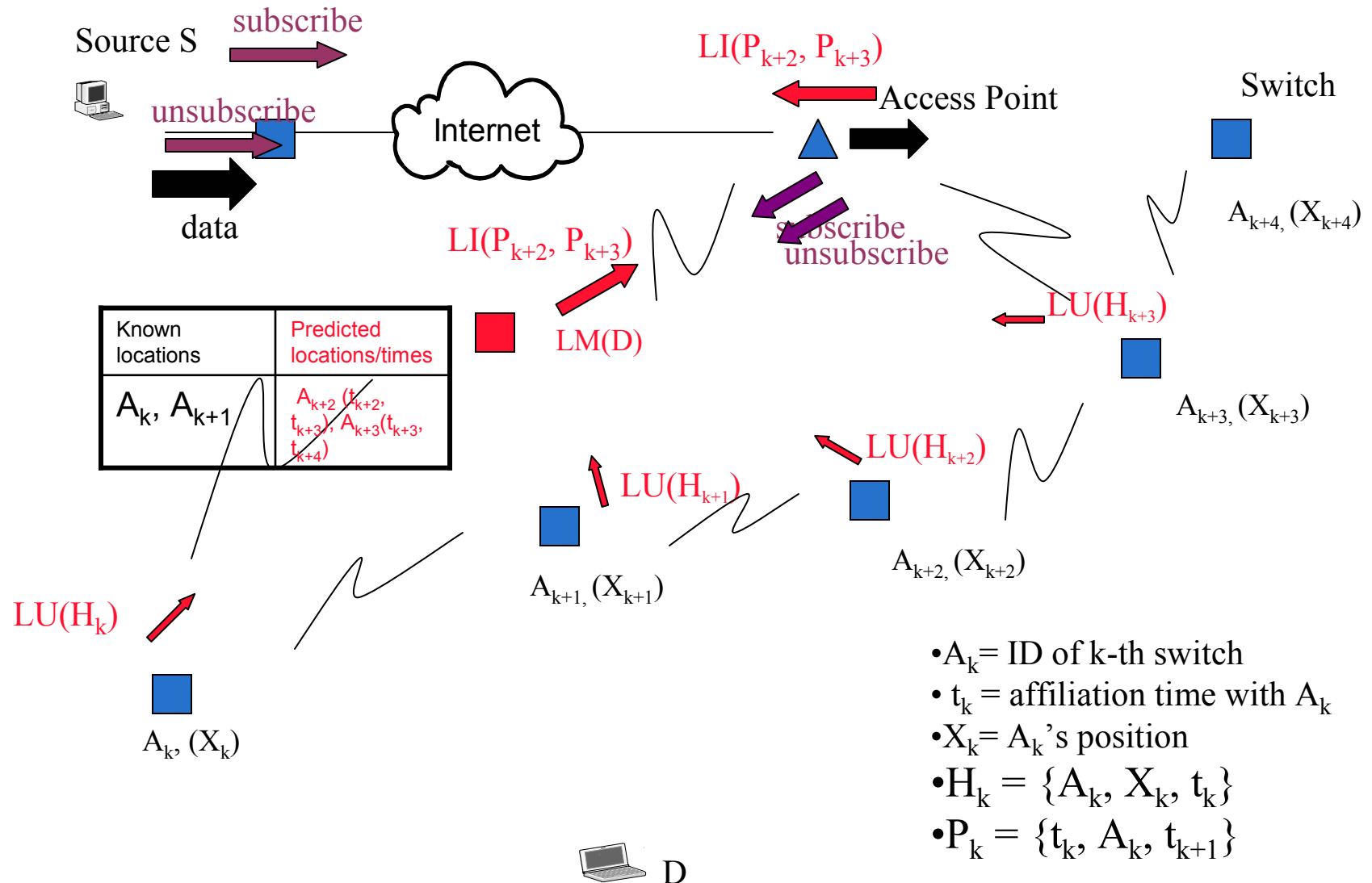
- *Spray Routing (SR)* attempts to address the problem of high mobility
  - SR multicasts data traffic in the vicinity of the last “known” location
  - Pros: SR improves the delivery rate
  - Cons: SR generates a significant overhead in duplicate data with increased mobility. SR is not effective beyond a certain effective mobility.
- A small amount of previous research has investigated ideas similar to “prediction of future locations based on history information”
  - *Niculescu, Nath* presents *Trajectory-Based Forwarding (TBF)*, a method to forwards packets in a dense MANET based on a predefined trajectory curve.
  - *Pathirana, et al* propose a *Robust Extended Kalman Filter (REKF)* as an estimator in predicting the mobile user’s trajectory.
  - *Liu, Maguire* propose a *pattern matching/recognition-based mobile motion prediction (MMP)* to estimate the user mobility.
- Anticipatory Routing is different from existing work in one key aspect: the reactability of the location tracking mechanism.

# Context For Our Study

- Studying “Anticipatory Routing” required provision of a context for “Anticipatory Routing” operation
  - Architecture is generic enough, i.e, not tied to a particular standard
- Our network consists of *Switches* and *Endpoints*
  - An endpoint affiliates with a one-hop away switch.
  - Switch can handle one or more endpoints
  - Endpoint’s affiliation is done proactively.
  - Switches are GPS-capable: GPS provides time synchronization and position information
- Traffic is stream-oriented (packetized voice, video)
- Our location tracking mechanism is rudimentary and borrows ideas from several known approaches
  - K switches act as location managers (LM).
  - An endpoint is mapped into one LM based on a simple hash. Endpoint updates LM through Location Update.
  - At session setup time, source subscribes with LM of destination.
- Routing protocol is a flat link-state protocol among switches.



# Anticipatory Routing



# Anticipatory Routing (cont'd)

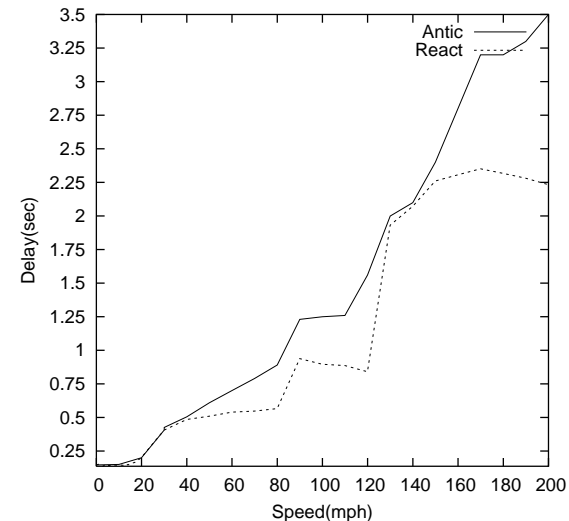
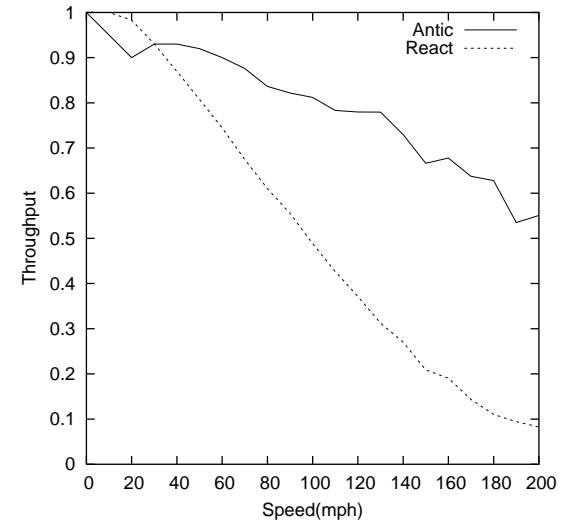
- Location Manager (LM) uses three basic steps to determine future locations and affiliation times
  - *Estimation of the endpoint's trajectory* as follows:
    - $L = 2$
    - 1. LM picks  $L$  and  $(L+1)$  latest positions  $(X_{k-L-1}..X_k)$  and  $(X_{k-L}..X_k)$
    - 2. For each set, LM derives two linear equations  $(m_1, h_1)$  &  $(m_2, h_2)$
    - 3. From  $(m_1, m_2)$ , directions  $(\theta_1, \theta_2)$  are determined
    - 4. If  $|\theta_1 - \theta_2| > \theta_{\text{thresh}}$ , then endpoint has “turned” and  $(m_1, h_1)$  becomes the trajectory equation. Otherwise,  $L=L+1$ ; goto step 1, etc..
  - *Estimation of the endpoint's direction  $\theta$* 
    - $\Theta$  is determined from slope  $m$ , current position  $X_k$  and oldest history position  $X_{k-t}$
  - *Estimation of future location, affiliation/departure time*
    - Future locations are derived from  $(m, h)$ , last known position  $X_k$  and network topology
    - Affiliation/departure times to/from switch regions are a function of inter-switch distance, estimated speed, switch diameter, etc..
- Upon data forwarding, switch-source selects final location based on delay estimate and set of anticipated locations and affiliation/departure times

# Anticipatory Routing (cont'd)

- Experimental Environment
  - Simulator built from scratch using C++ to independently control performance between fidelity and running time.
  - Multihop wireless network generated by placing a set of S switches and E endpoints in an L x L square area.
    - S = 100; E=40; L=1km; switch radius = 80 meters
  - Switches are static
  - Free-space propagation with threshold cutoff.
  - Mobility model is an extended random walk model.
  - Area model is a “wraparound” model: no area boundary.
  - Queuing delay at each node is modeled: packet is delayed by a time roughly proportional to the congestion.
  - Traffic is stream-oriented
    - Packets within a session generated by bernouilli trials
    - Data rate = 200 kbps
  - This parameters reflect future trends, including an increased reliability on wireless relays to extend the reach.

# Anticipatory Routing (cont'd)

- *Throughput (packet delivery ratio), Anticipated Fraction and Delay* are metrics of interest
- Analysis
  - AR does worse than traditional routing (TR) for speeds < 30mph
  - AR throughput ~92% around 30mph
  - AR performs better than TR above 30mph (throughput > 50%)
    - Loss-tolerant applications can benefit from this
  - Both AR and TR give rise to higher delays with increasing speed
    - Control message latency increases with frequency of affiliations, resulting in higher queueing delays



# Conclusion

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- Anticipatory Routing (AR) is a novel mechanism that addresses the problem of routing to highly mobile endpoints
  - AR improves throughput by 56% compared to traditional routing mechanisms in all cases studied
- In practice, AR might be applicable to
  - wireless broadband networks (e.g: wireless-MANs/“mesh extension”) with a large and dense mobile user base
  - tactical networks with truck-mounted routers where situation awareness update of video/imagery data is critical

